

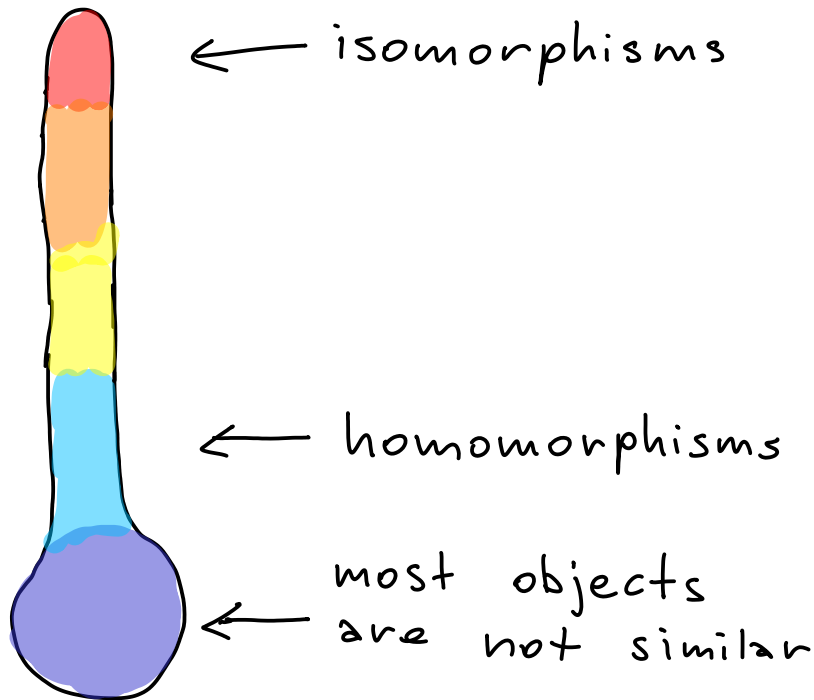
On the Complexity of  
Planar Regular Covering

Pavel Klavík

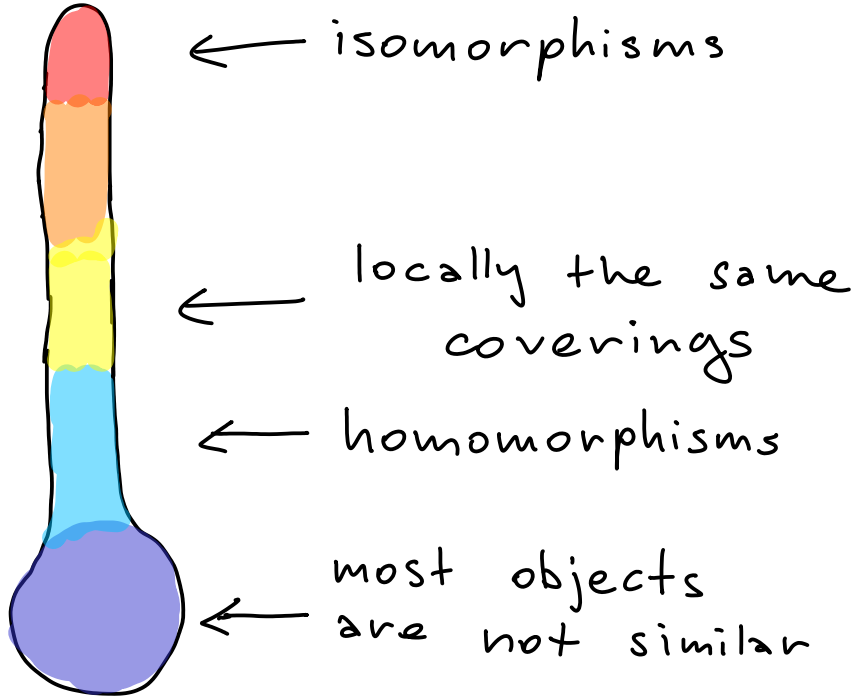
joint work with Jiří Fiala,  
Jan Kratochvíl and Roman Neděla.

ATCAGC 2013, Bovec

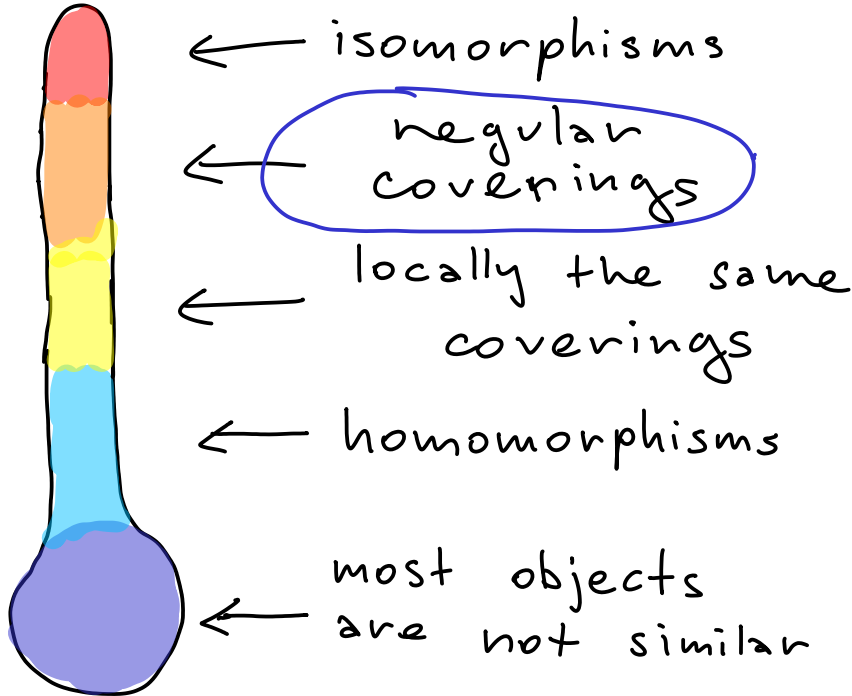
In mathematics, many papers study degree of similarity of objects.



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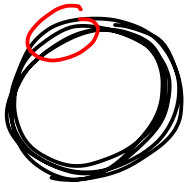


Notion of covering originates in topology.

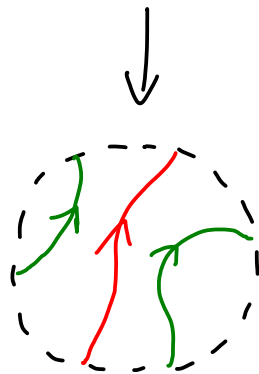
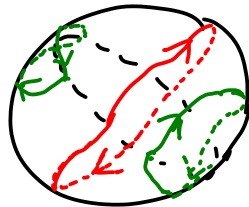
"Two surfaces are locally the same."



$$x \rightarrow (\cos x, \sin x)$$



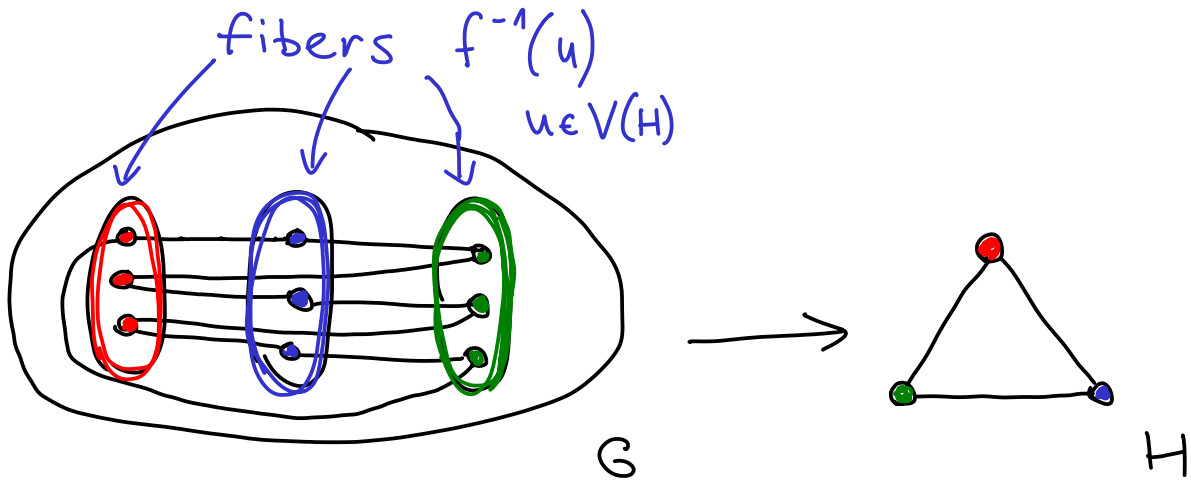
- real line covers a circle
- sphere covers the proj. plane



For graphs,  $G$  covers  $H$  if

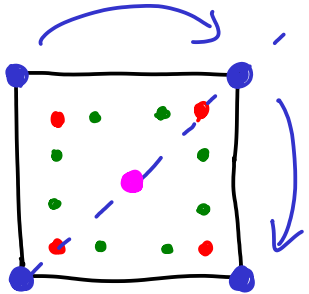
there exists a locally bijective  
homomorphism  $f: V(G) \rightarrow V(H)$

s.t.  $f|_{N[u]}$  maps bijectively to  $N[f(u)]$ .



# Group actions

- A group  $G$  transforms a set  $S$ .



$S$  is square

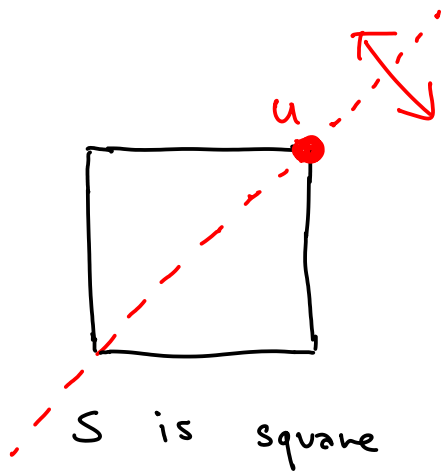
$G$  is  $D_4$

- Orbit of  $x \in S$  consists of all images of  $x$  under action of  $G$ .

-  $G$  acts transitively if it has exactly one orbit

# Group actions

- A group  $G$  transforms a set  $S$ .



$S$  is square

$G$  is  $D_4$

- stabilizer  $S(u)$  is a set of actions

s.t.  $u = u \circ g$

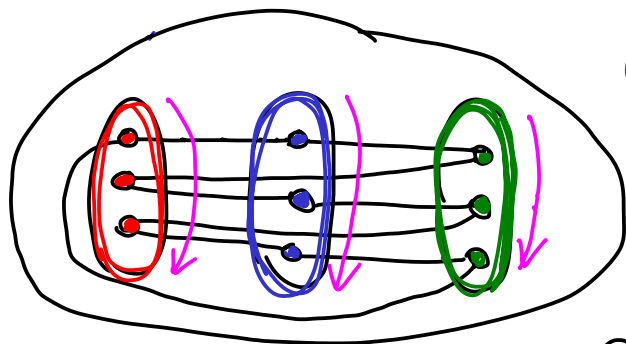
- group acts semiregul. if for  $\forall u \in S$  it has only the trivial stabilizer.



# Covering transformation group

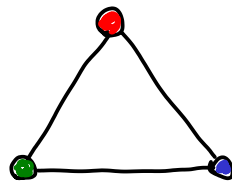
- covering transformations are fiber-preserving automorphisms

- they form a semiregular subgroup  $CT(G) \leq Aut(G)$



$G$

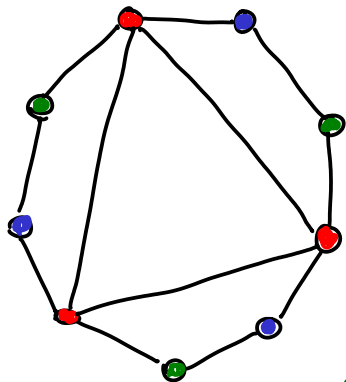
$$|CT(G)| \leq |fiber|$$



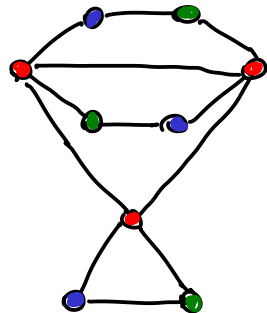
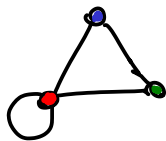
$H$

So what is a regular covering?

Covering is regular if CT acts transitively on every fiber,  
in other words  $|CT(G)| = |\text{fiber}|$ .



regular ✓

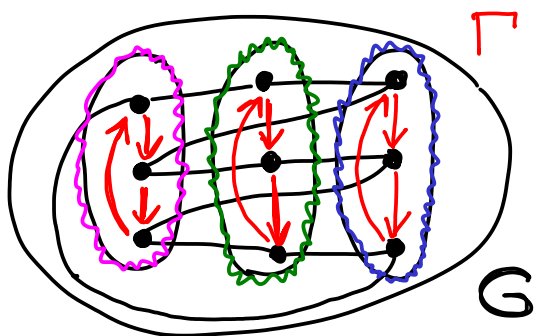


✗ non-regular

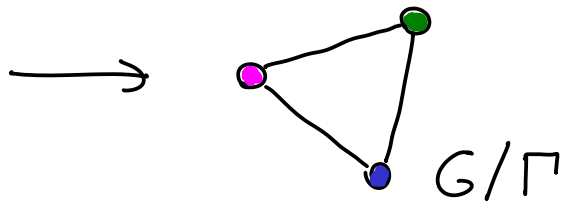
Definition using factorization:

Let  $\Gamma$  be a semiregular subgroup of  $\text{Aut}(G)$ . Then the factorization

$G/\Gamma$  defines a regular covering projection.

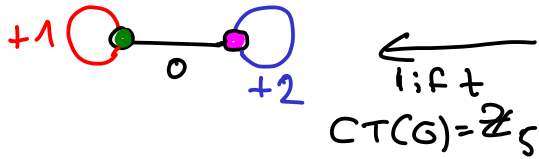


- orbits of  $\Gamma$  are fibers of the covering

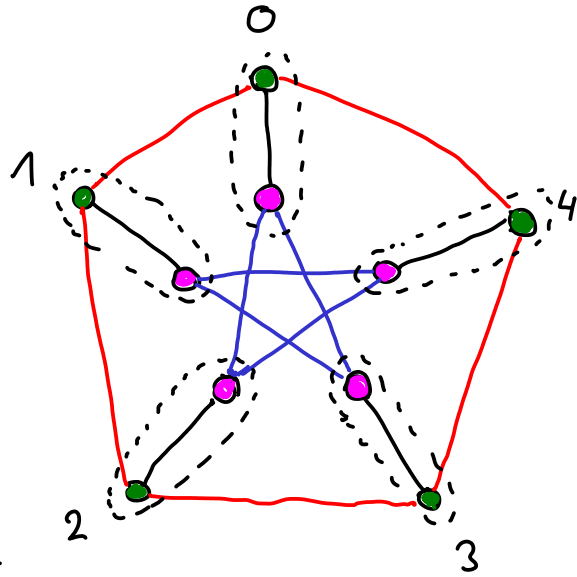


# Definition using voltages:

Based on unique walk lifting property.



$$CT(G) = \mathbb{Z}_5$$



- make  $|CT(G)|$  copies of the spanning tree of  $H$

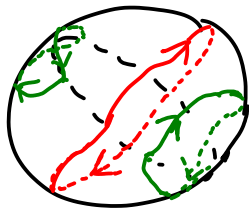
- put remaining edges

according to some element of  $CT(G)$

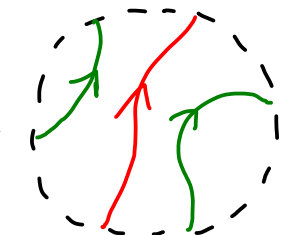
What about planarity?

Let the big graph  $G$  be planar.

Closely related to Negami's Conj:



sphere  
double-covers  
projective  
plane



A connected graph  $H$   
has a finite planar  
cover  $\iff H$   
can be embedded  
into the projective  
plane.

General case is still open:

- property "H has a planar cover" is closed under minor relation and  $Y\Delta$ -transformations
- only one forbidden minor  $K_{1,2,2,2}$  open.

Negami proved restricted versions:

- for 2-fold covers
- for regular covers

# Complexity problem H-COVER

Input: A graph  $G$ .

Question: Does  $G$  covers  $H$ .

There are some results:

Kratochvíl, Proskurowski, Telle, Fiala

- $k$ -reg  $H$  is NP-c for  $k \geq 3$
- dichotomy for small graphs

# The H-PLANAR COVER problem:

Input: A planar graph  $G$ .

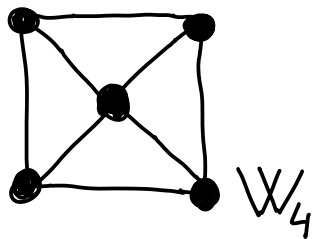
Question: Does  $G$  covers  $H$ .

The problem is trivial for  $H=K_2, \dots$

Non-trivial cases seems equally  
hard.

Bílka, Jirásek, K.,  
Tancer, Volec. WG'11

Small  
Open  
case:





Idea from last ATCAGC:

The PLANAR REGULAR COVER problem

Input: Graphs  $G$  and  $H$ ,  
 $G$  planar,  $H$  projective.

Question: Does  $G$  regularly covers  $H$ ?

We believe that this problem is solvable in polynomial time.

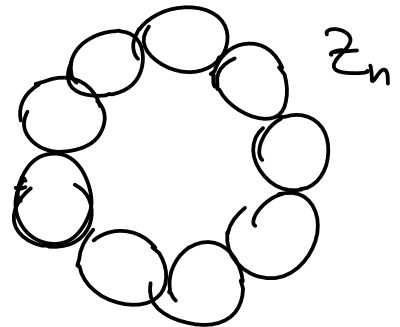
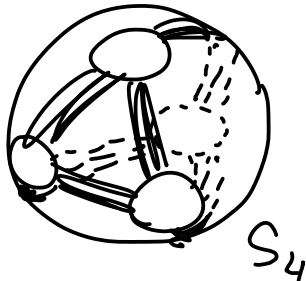
Let  $\mathcal{M}$  be an embedding of  $G$ ,  
then  $\text{Aut}(G) \cong \text{Aut}(\mathcal{M})$ .

-  $\text{Aut}(\mathcal{M})$  is very restricted:

- groups of platonic solids

$$S_4, \mathbb{Z}_2 \times A_5, \dots$$

-  $\mathbb{Z}_n, D_n$



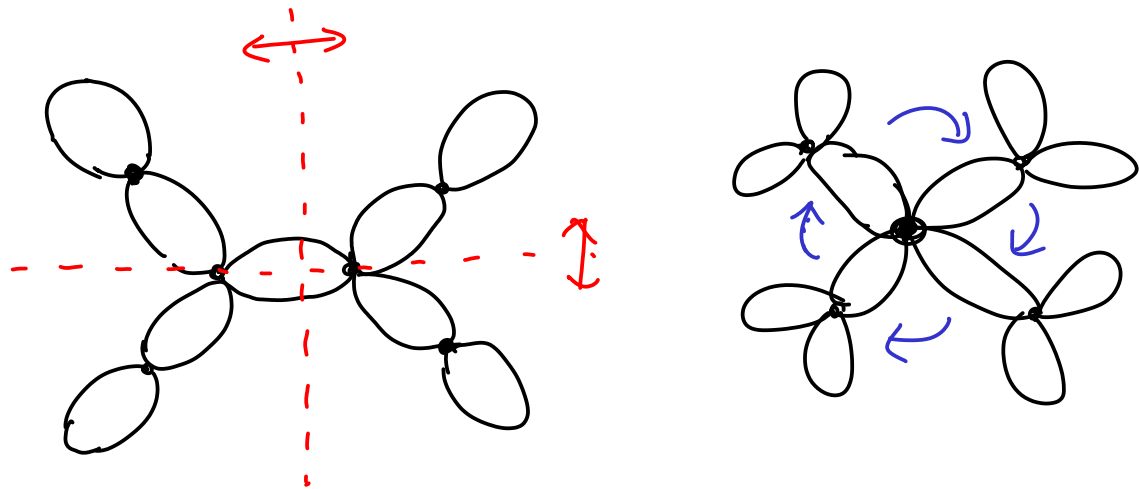
Idea is to compute  $\text{Aut}(G)$  and  
to guess  $\text{CT}(G)$ .

If  $G$  is 3-connected, the embedding  
is unique and  $\text{Aut}(G) = \text{Aut}(M)$ .

So we can guess  $\text{CT}(G)$  and to  
compute  $G/\text{CT}(G)$  and to  
test isomorphism with  $H$ .

1-connected pieces are easy:

Consider the block tree of  $G$ .



Every element of  $\text{Aut}(G)$  fixes central block or articulation.

Thank

You !

